

Week 2: Lab

COLLABORATION LEVEL 0 (NO RESTRICTIONS). OPEN NOTES.

1. Prove or disprove: $f = O(g)$ implies that $g = O(f)$.

What we expect: If true, a justification that holds for any functions f, g and for all $n > n_0$.
If false, a counterexample.

2. Prove or disprove: $f = O(g)$ implies that $g = \Omega(f)$.

What we expect: If true, a justification that holds for any functions f, g and for all $n > n_0$.
If false, a counterexample.

3. Find the order of growth of the following functions. For each function, give its $\Theta()$.

(a) $n \lg \lg n + n \lg n + \sqrt{n} \lg^2 n$

(b) $\sqrt{n} \lg n + n$

(c) $n^2 + \sqrt{n} \lg^3 n$

(d) $3^{\lg n} + n^2 + n \lg n$

(e) $\sqrt{3}^{\lg n} + n^2 + n \lg n$

(f) $2^n + 2^{2n}$

(g) $2^{\lg n} + \lg n^2$

(h) $(\lg n)^{\lg n} + n^3$

4. Arrange the following functions in ascending order of growth rate. For each pair of consecutive functions, give a brief justification on why they are in this order (For e.g., if you ordered A, B, C , you need to justify that $A = O(B)$ and $B = O(C)$).

$$2\sqrt{\log n}, 2^n, n^{4/3}, n(\log n)^3, n^{\log n}, 2^{2^n}, 2^{n^2}$$

5. Express the worst case running time of Insertion Sort as a sum and find its asymptotic order of growth (using the formula discussed this week).
6. Consider the **element uniqueness problem**: check whether all the elements in a given array are distinct. The problem can be solved by the following straightforward algorithm:

```

UNIQUEELEMENTS(A[0..n - 1])
1 // Checks whether all the elements in a given array are distinct
2 // Output: Returns “true” if all elements are distinct, and “false” otherwise
3 For i = 0 to n - 2
4     For j = i + 1 to n - 1
5         if A[i] == A[j] return false
6 return true

```

What is the running time (best-case, worst-case) of the algorithm, as a function of the n ? (hint: express as a sum and use the formula discussed).

7. Given two n -by- n matrices¹ A and B , by definition, their product $C = AB$ is an n -by- n matrix C where

$$C[i, j] = A[i, 0]B[0, j] + A[i, 1]B[1, j] + \dots + A[i, k]B[k, j] + \dots + A[i, n - 1]B[n - 1, j]$$

The straightforward algorithm to compute the product is given below:

```

MATRIXMULTIPLICATION(A[0..n - 1, 0..n - 1], B[0..n - 1, 0..n - 1])
1 // Multiplies two square matrices of order n by the definition-based algorithm
2 // Output: Matrix C = AB
3 For i = 0 to n - 1
4     For j = 0 to n - 1
5         // computes C[i,j]
6         C[i,j] = 0
7         For k = 0 to n - 1
8             C[i, j] += A[i, k] × B[k, j]

```

What is the running time (best-case, worst-case) of the algorithm, as a function of the order of the matrix, n ?

8. The following algorithm finds the number of binary digits in the binary representation of a positive decimal integer.

¹If you haven't seen matrices before, you can assume a matrix is a 2-dimensional array.

```

BINARY( $n$ )
1 // Input: A positive decimal integer  $n$ 
2 // Output: The number of binary digits in  $n$ 's binary representation
3 count = 1
4 While  $n > 1$ 
5     count++
6      $n = n/2$ 
7 return count

```

What is the running time (best-case, worst-case) of the algorithm, as a function of n ?

9. Consider the following algorithm:

```

MYSTERY( $n$ )
1 // Input: A positive decimal integer  $n$ 
2  $s = 0$ 
3 For  $i = 1$  to  $n$ 
4      $s = s + i \times i$ 
5 return  $s$ 

```

- (a) What does this algorithm compute?
 (b) What is the running time of the algorithm, as a function of n ?
10. Assume that we have an algorithm whose running time is $f(n)$ microseconds. Determine the largest size of a problem that can be solved by the algorithm in: (a) 1 second; (b) 1 hour; (c) 1 month.
 (a) $f(n) = \lg n$ (b) $f(n) = n^{10}$ (c) $f(n) = 2^n$.

Optional

11. You are presented with 9 marbles. All of the marbles look identical i.e. same shape, color, and dimensions (except for weight). However, 8 of the 9 marbles have exactly the same weight; the last marble is heavier. The only tool you have to measure weights is an old fashioned balance scale. You are only allowed to use the scale 2 times. How do you find the one marble that is not the same weight as the others?
 Generalize to n marbles: find it using the scale $\log_3 n$ times. (note: interview question)
12. You are given a set of n points on a circle in the plane. Come up with an algorithm that determines if there exists a pair of points that are antipodal (two points are antipodal if they are diametrically opposite). Analyze its running time (note: interview question) .